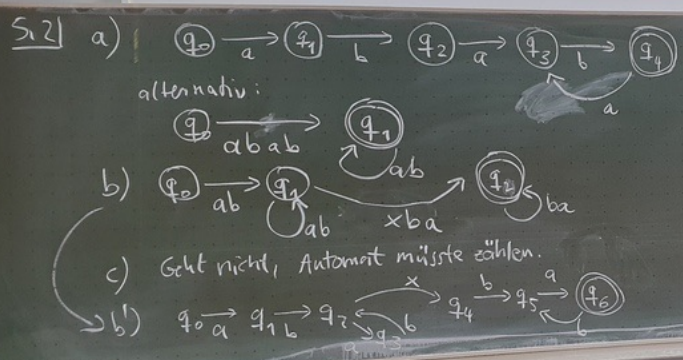


5.1] $L(A_1) = \{ a^n \mid n \in \mathbb{N}_0 \} = L(a^*)$
 $L(A_2) = \{ a w \mid w \in \{a, b\}^* \}$
 $\cup \{ b b^m a w \mid n \in \mathbb{N}_0, w \in \{a, b\}^* \}$
 $L(A_3) = \emptyset$
 $L(A_4) = \{ x a^n c b^m \mid x \in \{a, b, c\}, m, n \in \mathbb{N}_0 \}$

$= L(a(alb)^*)$
 $\cup L(b^+ a(alb)^*)$
 $= L((\epsilon/b^+) a(alb)^*)$
 $= L((alb)c(aa)^*cb^*)$

... die Sprache $L_2 = \{(ab)^n x (ba)^m \mid n, m \in \mathbb{N}\}$ erkennt. (Zur Erinnerung: $0 \notin \mathbb{N}$, denn $\mathbb{N} = \{1, 2, 3, \dots\}$ und $\mathbb{N}_0 = \{0, 1, 2, \dots\}$.)
 3. Welches Problem entsteht, wenn Sie bei der obigen Aufgabe verlangen, dass immer $n = m$ gilt, also für die Sprache $L_3 = \{(ab)^n x (ba)^n \mid n \in \mathbb{N}\} = \{abxba, ababxbaba, abababxbababa, \dots\}$?



Formal: $A_2 = (\Sigma, S, \delta, s_0, F)$
 mit $\Sigma = \{a, b, x\}$, $S = \{q_0, q_1, q_2, q_3, q_4, q_5, q_6\}$, $F = \{q_6\}$
 $s_0 = q_0$ und δ :

δ	q_0	q_1	q_2	q_3	q_4	q_5	q_6
a	q_1	-	q_3	-	-	q_5	-
b	-	q_2	-	q_4	q_5	-	q_6
x	-	-	q_4	-	-	-	-

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$S = \text{Pot}(\{s_0, s_1, s_2, s_3\})$
 $= \{ \emptyset, \{s_0\}, \{s_1\}, \{s_2\}, \{s_3\}, \{s_0, s_1\}, \{s_0, s_2\}, \{s_0, s_3\}, \{s_1, s_2\}, \{s_1, s_3\}, \{s_2, s_3\}, \{s_0, s_1, s_2\}, \{s_0, s_1, s_3\}, \{s_0, s_2, s_3\}, \{s_1, s_2, s_3\}, \{s_0, s_1, s_2, s_3\} \}$

$2^4 = 16$
 $\text{Pot}(Q) = 2^{|Q|}$

5.3] $\delta^*(s, \epsilon) = s$
 $\delta^*(s, aw) = \delta^*(\delta(s, a), w)$
 δ^* ist nur eine andere Schreibweise
 Vgl. $\sum_{i=1}^n a_i = \begin{cases} 0, & n=0 \\ a_{n-1} + \dots + a_1, & n \geq 1 \end{cases}$
 $\sum_{i=1}^n a_i = a_1 + a_2 + \dots + a_{n-1} + a_n, n \geq 1$

